

# When the Student becomes the Teacher

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- A report on our experience of teaching formal methods and an analysis of the exam results from the 2018–2019 academic year
  
- Observations and hypotheses to be further investigated.

## Module Overview

- Compulsory for third-year Computer Science BSc students.
- Optional for General Science BSc students.
- Prerequisites: Java programming module and Discrete Structures module.
- Runs over 12 weeks with 2 lecture hours and 2 lab hours per week.

- Design by Contract (1 week).
- Natural Deduction Proofs and the Coq theorem prover (3 weeks).
- Hoare Logic (2 weeks).
- Spec# (2 weeks).
- SAT/SMT (2 weeks).
- Model Checking (2 weeks).

Upon completion of the module, students should be able to:

- Explain the role of verification in software engineering.
- Create mathematically precise specifications.
- Prove the correctness of programs using Hoare Logic.
- Use different tools to analyse and verify properties of specifications.

# Assessment

- 30% continuous assessment (CA) and 70% best 3 out of 4 exam questions.
- CA: attend a 2 hour lab session each week to complete weekly assignment which is graded by the lab demonstrators.
- CA: 11 labs in total.
  - Lab 1–3: Natural Deduction Proofs with Coq
  - Lab 4–5: Hoare Logic
  - Lab 6–11: examine a range of verification tools including Spec# and Z3.

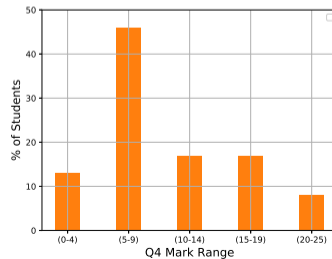
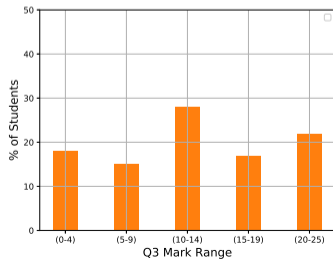
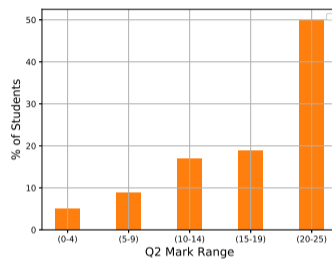
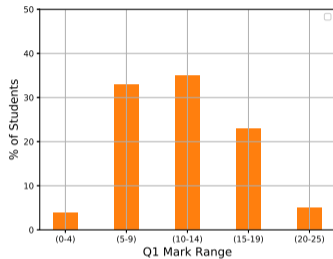


# Exam Structure

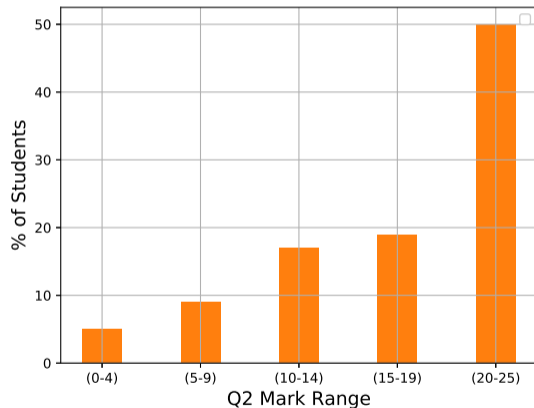
The pen and paper exam is 2 hours long and has the following structure:

Question	Examined Topics	Weight (marks)
Q1	Design by Contract Propositional and Predicate Logic Natural Deduction Proofs	25
Q2	Satisfiability CNF Translation DPLL (Pure literal, Unit clause and Unit propagation)	25
Q3	Hoare Logic	25
Q4	Basic SMT Encoding Spec# Programming (Pre/Post conditions, Loop invariants) Linear Temporal Logic Encoding	25

# Exam Results (2018–2019)

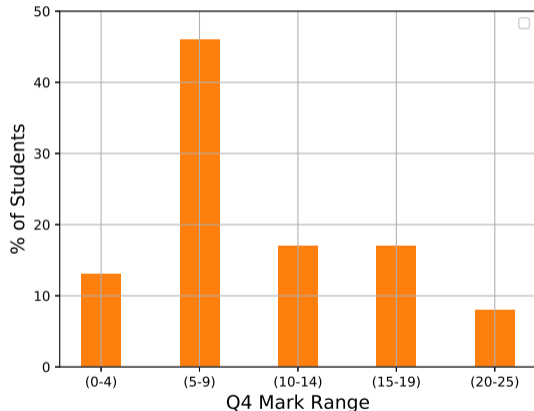


# Exam Results (2018–2019): Q2



- Students performed best on Q2 (satisfiability, CNF translation and DPLL).
- Q2 was the most popular question (88/92).
- Q2 is mechanical in nature (application of rules to formulae).

# Exam Results (2018–2019): Q4



- Students performed worst on Q4 (SMT encoding, Spec#, LTL model-checking).
- Q4 was the least popular question (24/92).
- Q4 was designed to challenge students on the topics:
  - Encode specification into SMT formulas.
  - Write Spec# program with specifications corresponding to C# sum of integer array.
  - SAT encodings for reachability, safety and liveness properties.

# Reflecting on Teaching and Learning

## Learning:

- Module is generally challenging for students:
  - identifying loop invariants
  - presenting Hoare Logic proofs
  - understanding low-level SAT/SMT encodings
- Verification tools are not very reliable and don't always give useful feedback.
- Practical applications of formal methods not clear to students.

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## Teaching:

- Difficult to help students see the value in formal methods since they are not widely used in industry.
- Tools are normally not scalable for real-world examples making it difficult to demonstrate their usefulness to students.
- Necessary to combine slides and worked-through whiteboard examples to explain detailed computation steps.

# Whiteboard Example

$x_0 \in D$   
 $t: D, x=P$

$\forall x (P(x) \rightarrow Q(x)), \exists x (P(x)) \vdash \exists x (Q(x))$

- $\forall x (P(x) \rightarrow Q(x))$  Premise
- $\exists x (P(x))$  Premise  $\rightarrow \exists E$
- $x_0, P(x_0)$  assumption
- $P(x_0) \rightarrow Q(x_0)$  H.E. 1.
- $Q(x_0)$   $\Rightarrow$  C. 3, 4.
- $\exists x (Q(x))$   $\exists I$  3, 5.
- $\exists x (Q(x))$   $\exists E$  2, 3-6.
- Q.E.D.

$(A \wedge B) \rightarrow C$   
 $\neg(A \wedge B) \vee C$   
 $\neg A \vee \neg B \vee C$   
 $\neg A \vee (B \rightarrow C)$   
 $A \rightarrow (B \rightarrow C)$

$C = \exists x (Q(x))$



## Observation 1:

- Automated verification tools are not appealing to students.
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## Observation 2:

- Students good at natural deduction but not using Coq.
- Difficult to connect pen and paper proof with Coq proofs.

## Observation 3:

- We did not integrate our own research tools into this module.
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## Observation 4:

- Mixed reactions to Hoare Logic.
- Used whiteboard examples.
- Students with strong mathematical background performed better on this topic.

# Whiteboard Example

① before loop  $\{i = \frac{x!}{a^i}\} \Rightarrow \{i=1\} = \text{True}$ .

$a = x;$

$\{i = \frac{x!}{a^i}\}$

$y = 1;$

$\{y = \frac{x!}{a^i}\}$

② in the loop  $\{a > 0\}$

$\{y * a = \frac{x!}{(a-1)!} \wedge a > 0\}$

$y * a;$

$\{y = \frac{x!}{(a-1)!}\}$

$a--;$

$\{y = \frac{x!}{a^i}\}$

}

③ after loop:  $\{y = \frac{x!}{a^i} \wedge a \leq 0\}$   
 $\downarrow$   
 $\{y = x!\}$

① for loop  
 $x = i;$   
 $\{x = A^0\}$   
 $i = 0;$   
 $\{x = A^i\}$

② while ( $i \neq N$ )  
 $\{x * A = A^{i+1} \wedge i! = N!\}$   
 $x * A; x = x * A;$

$\{x = A^{i+1}\}$

$i++;$

$\{x = A^i\}$

③  $\{x = A^i \wedge i = N\}$

$\downarrow$   
 $\{x = A^N\}$

$x = A^i$   
 $x * A = A^i * A$   
 $x * A = A^{i+1}$

$x = A^i$

	x	i	A	N
0	1	0	3	5
1	3	1	3	5
2	9	2	3	5
3	27	3	3	5
4	81	4	3	5
5	243	5	3	5

- Observations 1 and 2 point to a lack of tool usability. This can also hinder the uptake of formal methods in industry.
- Observation 3 has caused us to integrate our research tool, MaxUSE, into teaching. So far, students seem interested in this.
- Observation 4 noted that students found Hoare Logic difficult but this may have been exacerbated by their difficulty with some of the tools e.g. Spec#.

We intend to investigate the following hypotheses in the current (2019–2020) academic year:

## Hypothesis 1:

- The development of an online repository of real-world examples would be useful for both teaching and illustrating industrial uses of formal methods.
- Action: design and distribute a survey among past students to identify the most interesting and educational examples to be used in class.

## Hypothesis 2:

- A platform that turns logical reasoning proofs into games would increase the interactions between students and lecturers.
- Action: a live coding session using SMT solvers to solve a Sudoku.

## Conclusions and Future Work

- We discussed our own experience of both studying and teaching the same software verification module at Maynooth University.
- Based on the 2018–2019 results, we made several observations from which we have constructed two hypotheses that we intend to investigate in the current (2019–2020) academic year.
- We plan to work with the education research group at Maynooth to develop interesting experiments in order to improve our teaching of this course.
- Maynooth offers a similar module at Master's level and we plan to investigate it and compare the results to this paper.



# Questions?